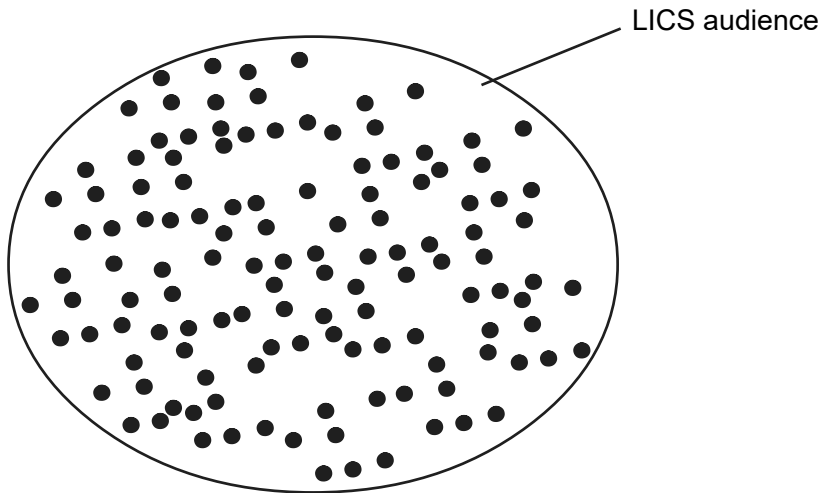

The Big-O Problem is Decidable for Max-Plus Automata (PSPACE-complete)

Laure Daviaud
University of East Anglia

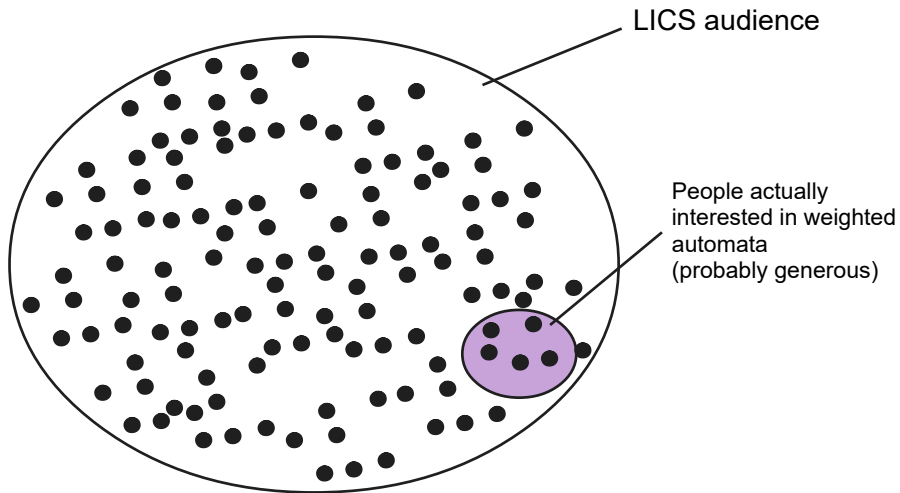
David Purser
University of Liverpool

In this room

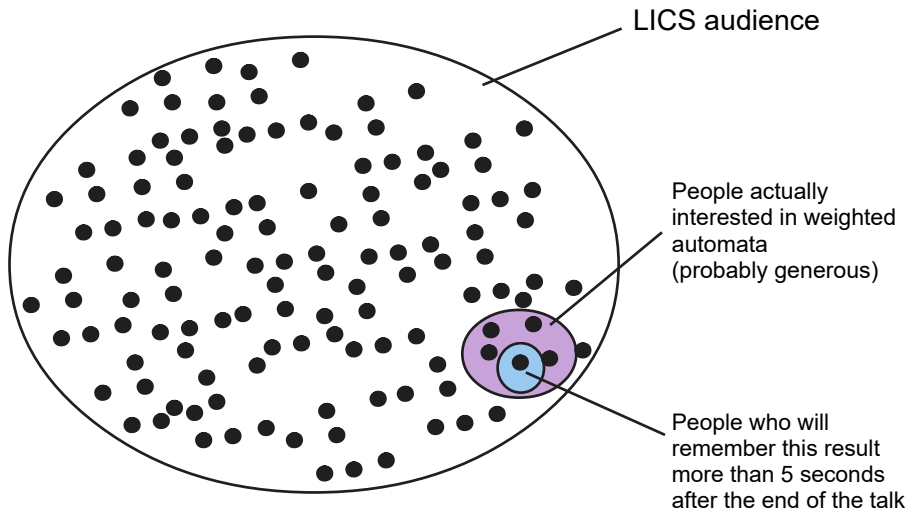
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Game plan

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(easier said than done...)

The property and the model

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Extension of Boolean automaton with non-negative integers on transitions, combined using operations max and sum.

$$\Sigma^* \rightarrow \mathbb{N} \cup \{-\infty\}$$

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There exists C such that for all $w \in \Sigma^*$, $f(w) \leq Cg(w) + C$

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- Decidable: approximation, min-plus
[Hashiguchi, Leung, Simon, Colcombet-D.-Zuleger]

One insight: turning infinite into finite

What happens when you repeat a word in an automaton?

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$(a^\# b)^\# = (\infty, \infty)$ but $(a^\# b)^\flat = (\infty, 1)$

→ New flattening operation \flat

The End (or just the beginning)

— Theorem [D., Purser]

The big-O problem is PSPACE-complete for max-plus automata.